Probabilistic Programming Hongseok Yang KAIST

Last part based on work with Chris Heunen, Ohad Kammar and Sam Staton

This Review ... discusses some of the state-ofthe-art advances in the field, namely, probabilistic programming, Bayesian optimization, data compression and automatic model discovery.

> Zoubin Gharahmani 2015 Nature Review

What is probabilistic programming?

(Bayesian) probabilistic modelling of data

- I. Develop a new probabilistic (generative) model.
- 2. Design an inference algorithm for the model.
- 3. Using the algo., fit the model to the data.

(Bayesian) probabilistic modelling of data in a prob. prog. language

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a generic inference algo. of the language

Line fitting











s b





Q: posterior of (s,b) given $y_0=0.6$, ..., $y_6=8.4$?

$$p(s, b | y_0, ..., y_6) = \frac{p(y_0, ..., y_6 | s, b) \times p(s, b)}{p(y_0, ..., y_6)}$$

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Anglican program

Anglican program

> (observe (normal (f 0) .5) .6) (observe (normal (f 1) .5) .7) (observe (normal (f 2) .5) 1.2) (observe (normal (f 3) .5) 3.2) (observe (normal (f 4) .5) 6.8) (observe (normal (f 5) .5) 8.2) (observe (normal (f 6) .5) 8.4)

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[s b])

Samples from prior



Samples from posterior



Underfit?



Underfit?



Underfit?



Expressive prob. PLs enable one to explore advanced models easily.
> (observe (normal (f 0) .5) .6) (observe (normal (f 1) .5) .7) (observe (normal (f 2) .5) 1.2) (observe (normal (f 3) .5) 3.2) (observe (normal (f 4) .5) 6.8) (observe (normal (f 5) .5) 8.2) (observe (normal (f 6) .5) 8.4)

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[s b]) Anglican fully supports higher-order functions
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Anglican fully supports higher-order functions

```
(let [F (fn []
      (let [s (sample (normal 0 2))
            b (sample (normal 0 6))]
        (fn [x] (+ (* s x) b)))
  f (F)]
(observe (normal (f 0) .5) .6)
(observe (normal (f 1) .5) .7)
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[Sb]) f)

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Samples from posterior



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Samples from posterior



Captcha breaking





Le, Baydin, Wood [2016]




Sample a string.





I. Sample a string.

2. Generate an image using complex JVM code.





- . Sample a string.
- 2. Generate an image using complex JVM code.

SMKBDF







Neural net as a part of inference engine.

Approximates the inverse of the Captcha program.





I. Sample a 3D object.



- I. Sample a 3D object.
- 2. Score the object.



- I. Sample a 3D object.
- Score the object.
 Used stochastic future.



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Why might QONFEST audience be interested?

Reason I: Unusual use of concurrency.

An inference algo. for a prob. PL can be viewed as a non-standard approximate interpreter.

Sequential Monte Carlo (in short SMC) is a popular algo. with many variants.

SMC runs sequential probabilistic programs concurrently with added synchronisation.



























Goal: Generate samples from posterior p(x,y | a=2.1,b=1.8).



w:1







(let [x (sample (normal 0 2))
 <u>a (observe (normal x 1) 2.1)</u>
 y (sample (normal (* 0.9 x) 2))
 b (observe (normal y 1) 1.8)]
 [x y])
Sequential
Monte-Carlo
algorithm

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I. Update weights. w \leftarrow pnormal(2.1;x,1)



I. Update weights. $w \leftarrow p_{normal}(2.1;x,1)$







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Sequential Monte-Carlo algorithm





Sequential Monte-Carlo algorithm












































Inference algorithms for Anglican

Method	Туре	Description
importance	IS	Importance sampling (likelihood weighting)
smc	IS	Sequential Monte Carlo
pcascade	IS	Particle cascade (asynchronous sequential Monte Carlo)
pgibbs	PMCMC	Particle Gibbs (iterated conditional SMC)
pimh	PMCMC	Particle independent Metropolis-Hastings
pgas	PMCMC	Particle Gibbs with ancestor sampling
ipmcmc	PMCMC	Interacting particle Markov chain Monte Carlo
lmh	MCMC	Lightweight Metropolis-Hastings
rmh	MCMC	Random-walk Lightweight Metropolis-Hastings
almh	МСМС	Adaptive scheduling lightweight Metropolis- Hastings
palmh	МСМС	Parallelised adaptive scheduling lightweight Metropolis-Hastings
plmh	MCMC	Parallelised lightweight Metropolis-Hastings
bamc	MAP	Bayesian Ascent Monte Carlo
siman	MAP	MAP estimation via simulated annealing

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rmh	MCMC	Random-walk Lightwe 6 variants of SMC
almh	МСМС	Adaptive scheduling lightweight Metropolis- Hastings
palmh	МСМС	Parallelised adaptive scheduling lightweight Metropolis-Hastings
plmh	MCMC	Parallelised lightweight Metropolis-Hastings
bamc	MAP	Bayesian Ascent Monte Carlo
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Key question

How to reason about prob. programs under such concurrent non-standard approximate semantics?

Concrete question I: Correctness

[Q] Are these algo. correct for prob. programs?
Usually proved for Rⁿ or simple cases.
Challenge I: Expressiveness of prob. PLs.
Challenge 2: Subtle notion of correctness.

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Challenge 1: Expressiveness of prob. PLs.
Challenge 2: Subtle notion of correctness.

Seq. Monte Carlo gives a right answer (weak convergence) as the # of threads goes to ∞ .

Programs may be semantically equivalent but some are easy for these algo., and some hard.

[Q] Capture this difference by an algo.-specific refinement \Box . Develop proof rules for \Box .

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 $Prog2 \sqsubseteq ProgI$

Concrete question 3: Good sublanguages

[QI] Find a sublanguage that drops the time complexity of an inference algorithm.

[Q2] Find a sublanguage that allows a GPUbased implementation of an inference algo.

Reason 2: Probabilistic PLs raise new semantic issues.

(let [F (fn [] (let [s (sample (normal 0 2)) b (sample (normal 0 6))] (fn [x] (+ (* s x) b))) f (add-change-points F 0 6)] (observe (normal (f 0) .5) .6)(observe (normal (f 1) .5) .7)(observe (normal (f 2) .5) 1.2)(observe (normal (f 3) .5) 3.2) (observe (normal (f 4) .5) 6.8)(observe (normal (f 5) .5) 8.2) (observe (normal (f 6) .5) 8.4)

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I. Higher-order functions.

Issue I: Higher-order functions

Measure theory provides a standard foundation of probability theory.

But it doesn't support HO fns well.

$$ev: (\mathbb{R} \rightarrow_m \mathbb{R}) \times \mathbb{R} \rightarrow \mathbb{R}, \quad ev(f,x) = f(x).$$

[Aumann 61] ev is not measurable no matter which σ -algebra is used for $\mathbb{R} \rightarrow_m \mathbb{R}$.

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I. Higher-order functions.2. Conditioning and prog. eqs.
Issue 2: Conditioning and prog. eqs [[e:rea1]] ∈ M(ℝ)

- M should model prob. computations.
- M should validate equations from statistics.
- M should be commutative.
- Difficult to find such M due to conditioning.

Issue 2: Conditioning and prog. eqs [[e:rea1]] ∈ M(ℝ)nonfinite measures

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Issue 2: Conditioning and prog. eqs nearly-finite measures $[e:real] \in M(\mathbb{R})$ nonfinite measures

- M should model prob. computations.
- M should validate equations from statistics
- M should be commutative.
- Difficult to find such M due to conditioning.

Issue 3:

Meta-programming features

My ML colleagues are very much interested in probabilistic models for programs.

They express such models using quote & eval.

[Q] How to interpret meta-programming features in Anglican/Church/Venture?

My research*: Address issues 1&2 with Quasi-Borel spaces.

* based on Heunen et al.'s LICS'17

Big picture 1: Extend measure theory using category theory.

Higher-order fns.
Conditioning, prog. eqs.

Higher-order fns.
Conditioning, prog. eqs.

Meas_B



























Big picture 2: Random element first.

Random element α in X

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- X set of values.
- Ω set of random seeds.
- Random seed generator.

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 $\alpha:\Omega\to X$

- X set of values.
- Ω set of random seeds.
- Random seed generator.

I. Σ⊆2Ω, Θ⊆2^X 2. μ : Σ→[0, I]

 $\alpha : \Omega \to X$ is a random element if $\alpha^{-1}(A) \in \Sigma$ for all $A \in \Theta$

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Random element α in X

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 $\alpha:\mathbb{R}\to X$

- X set of values.
- \mathbb{R} set of random seeds.
- Random seed generator.

I. \mathbb{R} as random source 2. Borel subsets $\mathfrak{B} \subseteq 2^{\mathbb{R}}$

 $\alpha: \mathbb{R} \to X$

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- \mathbb{R} set of random seeds.
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 $I \,.\, \mathbb{R}$ as random source

2. Borel subsets $\mathfrak{B} \subseteq 2^{\mathbb{R}}$

 $\alpha: \mathbb{R} \to X$

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I. \mathbb{R} as random source 2. Borel subsets $\mathfrak{B} \subseteq 2^{\mathbb{R}}$ 3. $\mathbb{M} \subseteq [\mathbb{R} \rightarrow X]$

 $\alpha : \mathbb{R} \to X \text{ is a random variable}$ if $\alpha \in M$

- X set of values.
- \mathbb{R} set of random seeds.
- Random seed generator.

I. \mathbb{R} as random source 2. Borel subsets $\mathfrak{B} \subseteq 2^{\mathbb{R}}$ 3. $\mathbb{M} \subseteq [\mathbb{R} \rightarrow \mathbb{X}]$

- Measure theory:
 - Measurable space (X, $\Theta \subseteq 2^X$).
 - Random element is an induced concept.
- QBS:
 - Quasi-Borel space (X, $M \subseteq [\mathbb{R} \rightarrow X]$).
 - M is the set of random elements.

Quasi-Borel spaces

- New axiomatisation of probability theory.
- Enabled us to generalise classical results in probability theory such as de Finetti thm.

Try probabilistic PLs

Anglican: http://www.robots.ox.ac.uk/~fwood/anglican/ WebPPL: http://webppl.org/