

Graphical Linear Algebra

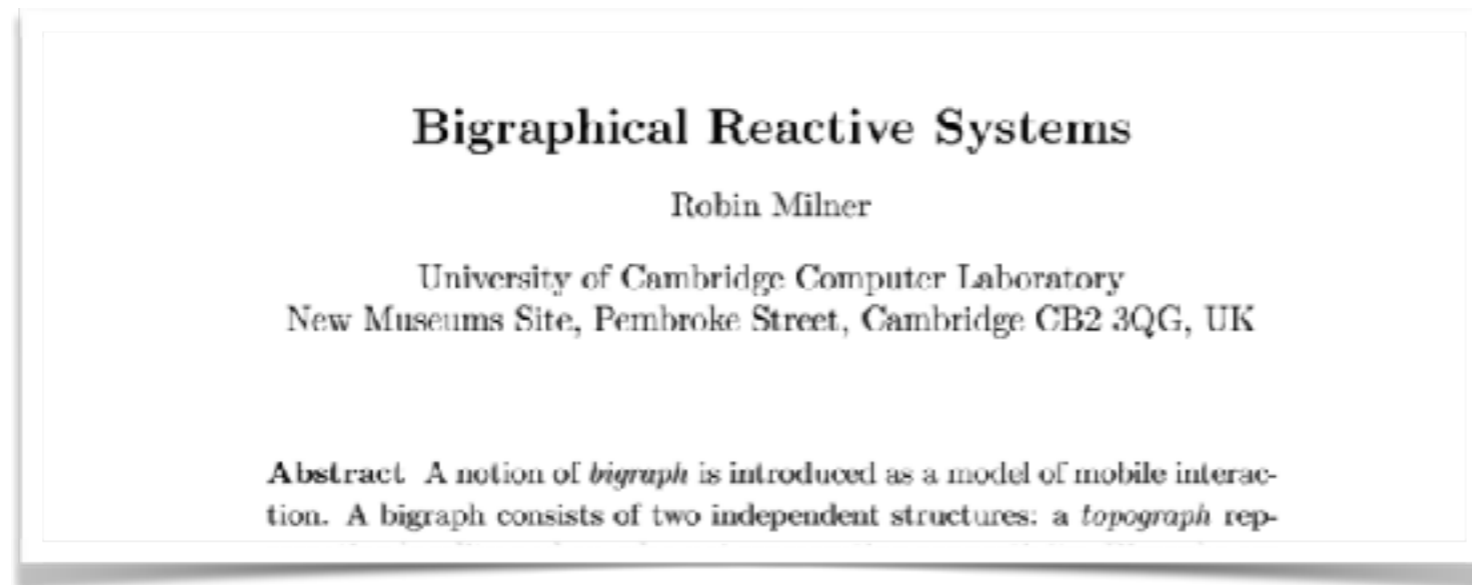
a specification language for linear algebra

Pawel Sobocinski

based on joint work with Filippo Bonchi and Fabio Zanasi

CONCUR

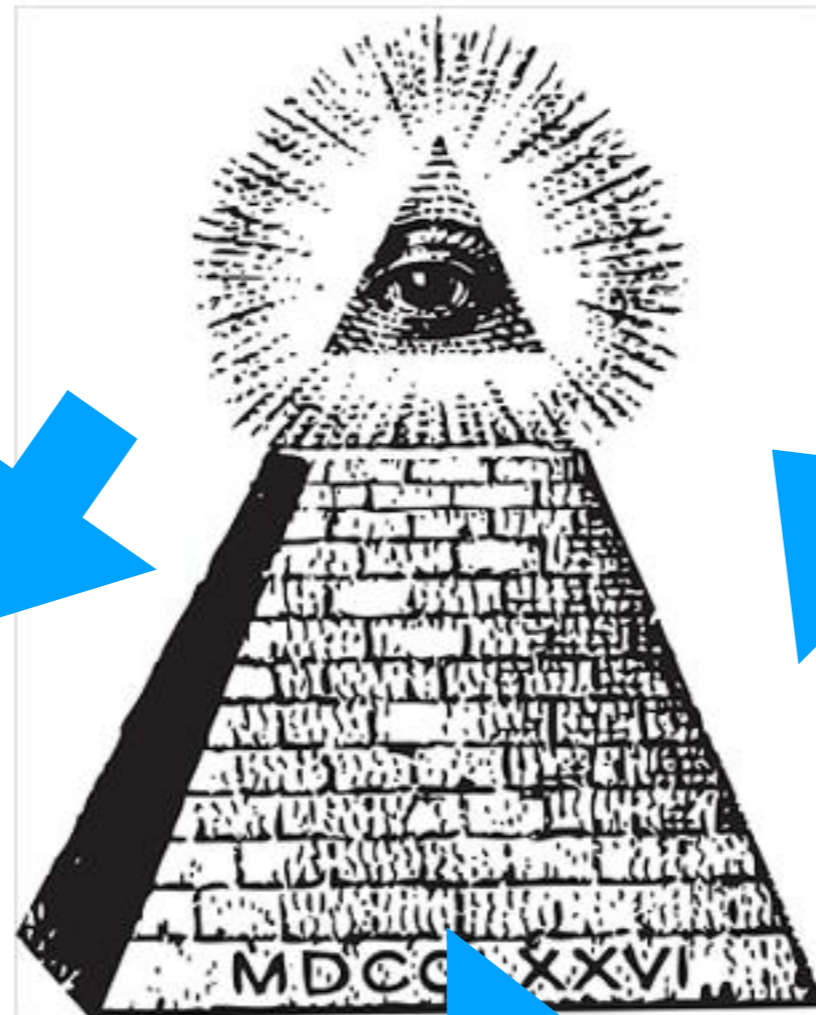
- My first CS conference: Concur 2001, Aalborg



- My first talk: EXPRESS 2002, Brno (a Concur 2002, satellite workshop)
- My first paper at Concur 2003, Marseille (with Bartek Klin)



Concur (process algebra)



string diagrams

linear algebra

(Jean-Raymond Abrial : how theorem provers make you treat mathematics as a branch of software engineering)

This talk: how to see linear algebra as a branch of process algebra

implementations \subseteq **specifications**

runnable, completely
specified

non-deterministic,
partially specified
behaviour

functions \subseteq **relations**

single-valued,
total

non single-valued,
non total

string diagrams

- showing up in a growing number of recent CS papers
 - Abramsky, Duncan, Coecke, ... - Categorical Quantum Foundations and Quantum Computation
 - Mellies, ... - Logic, Game Semantics
 - Ghica, Jung, ... - Digital Circuits
 - Baez, Bonchi, Erbele, Fong, S., Zanasi, ... - Signal Flow Graphs, Control and Systems Theory
 - Coecke, Sadrzadeh, ... - Computational Linguistics
 - ...



1st Workshop on String Diagrams in Computation Logic and Physics

Jericho Tavern, Oxford

8-9 September, 2017

(satellite of FSCD,
next year satellite of CSL)

linear algebra

- the most practical mathematical theory?
 - the engine room of systems and control theory, quantum computing, network theory, ...
 - mathematical physics and engineering relies on it: systems of nonlinear differential equations are solved with linear approximations
 - shows up in surprising places (Petri net invariants, PageRank is an eigenvector, SVD in data science and learning, ...)
- Graphical Linear Algebra - linear algebra with string diagrams
 - focus on linear **relations** rather than on linear **maps**
 - GraphicalLinearAlgebra.net

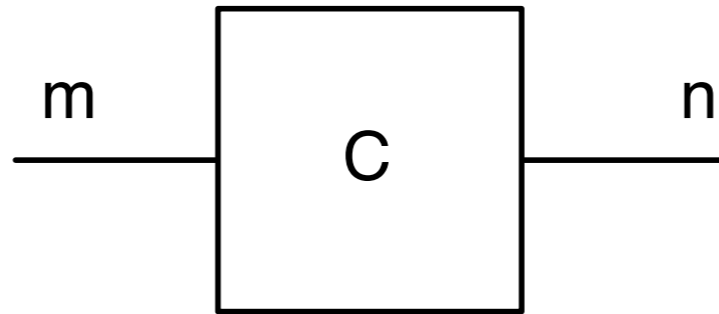
Plan

- String diagrams & diagrammatic reasoning
 - **what is it?**
 - why is it relevant for cs?
- Graphical Linear Algebra
- Fun stuff

props

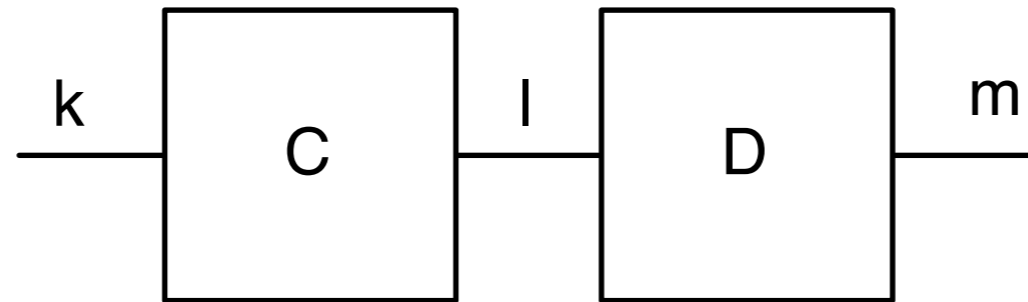
- A prop is a strict symmetric monoidal category with
 - strict means: \otimes is associative on the nose
 - objects = natural numbers
 - $m \otimes n := m + n$ (I will usually write $m \oplus n$)
- Simple examples:
 - permutations of finite sets
 - functions between finite sets
- prop homomorphism = identity on objects symmetric monoidal functor

A string diagram



Synchronising Composition

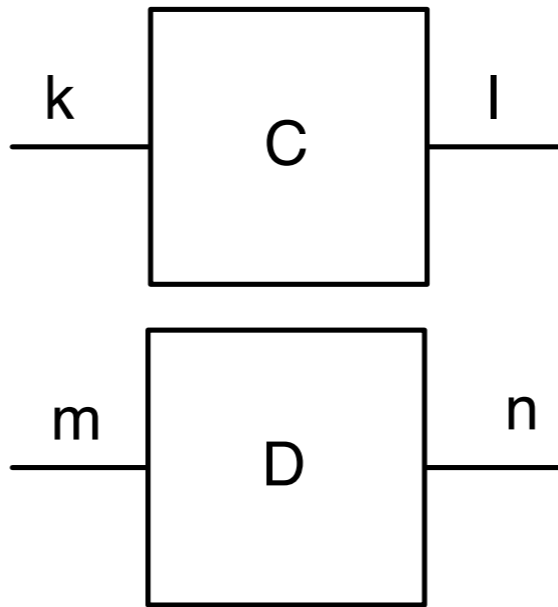
$$\frac{C : k \rightarrow l \quad D : l \rightarrow m}{C ; D : k \rightarrow m}$$



“C and D synchronise on l”

Parallel composition

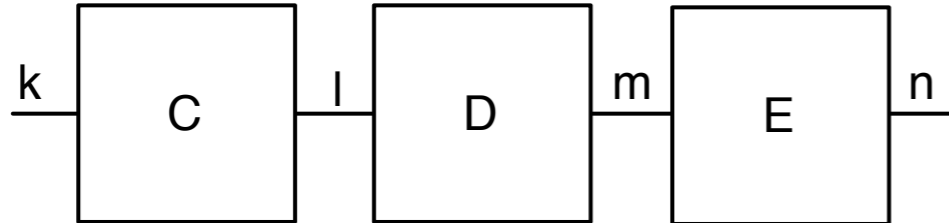
$$\frac{C : k \rightarrow l \quad D : m \rightarrow n}{C \oplus D : k \oplus m \rightarrow l \oplus n}$$



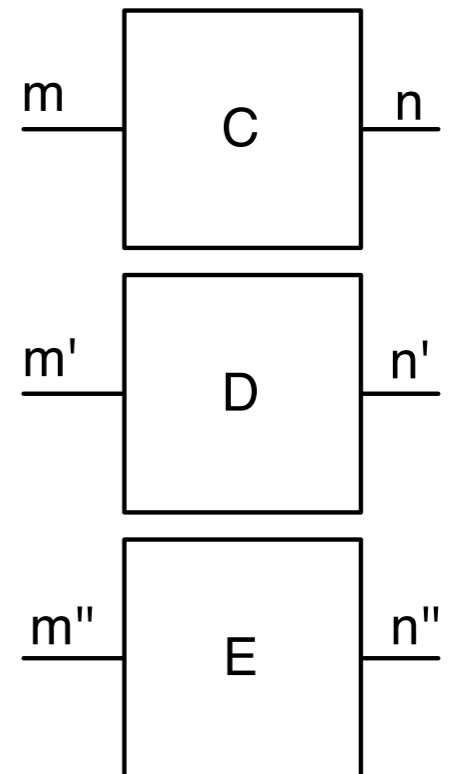
“C and D in parallel”

Perks of the notation

$$\frac{C : k \rightarrow l \quad D : l \rightarrow m \quad E : m \rightarrow n}{(C ; D) ; E = C ; (D ; E) : k \rightarrow n}$$

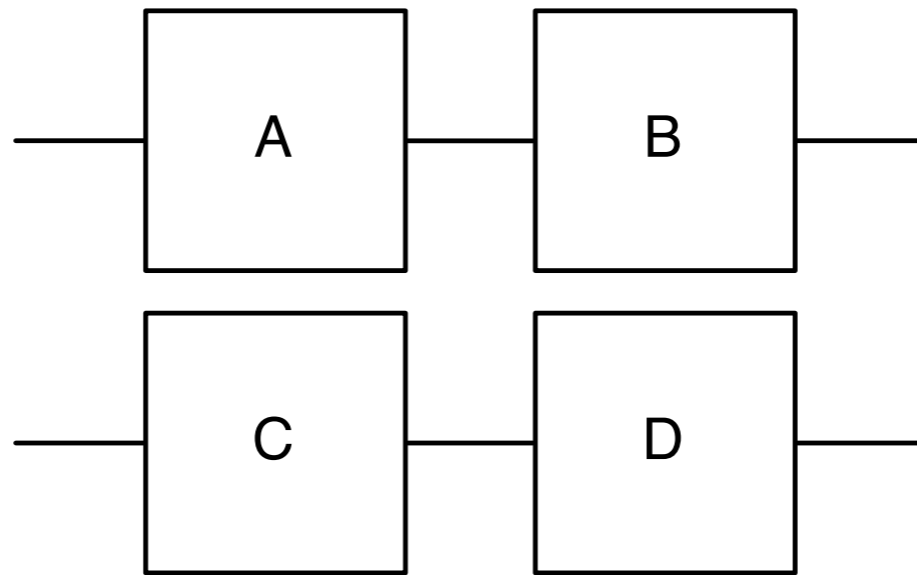


$$\frac{C : m \rightarrow n \quad D : m' \rightarrow n' \quad E : m'' \rightarrow n''}{(C \oplus D) \oplus E = C \oplus (D \oplus E) : m \oplus m' \oplus m'' \rightarrow n \oplus n' \oplus n''}$$



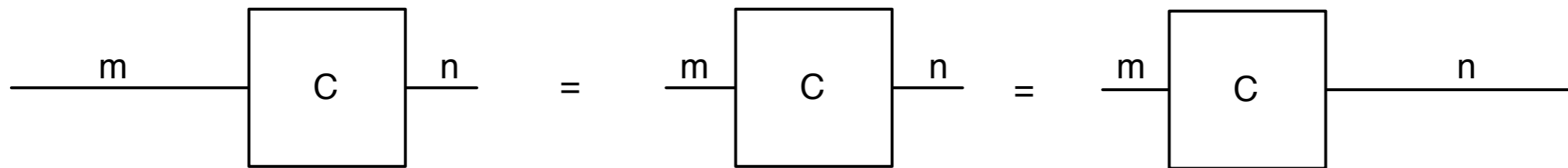
More perks

$$(A ; B) \oplus (C ; D) = (A \oplus C) ; (B \oplus D)$$

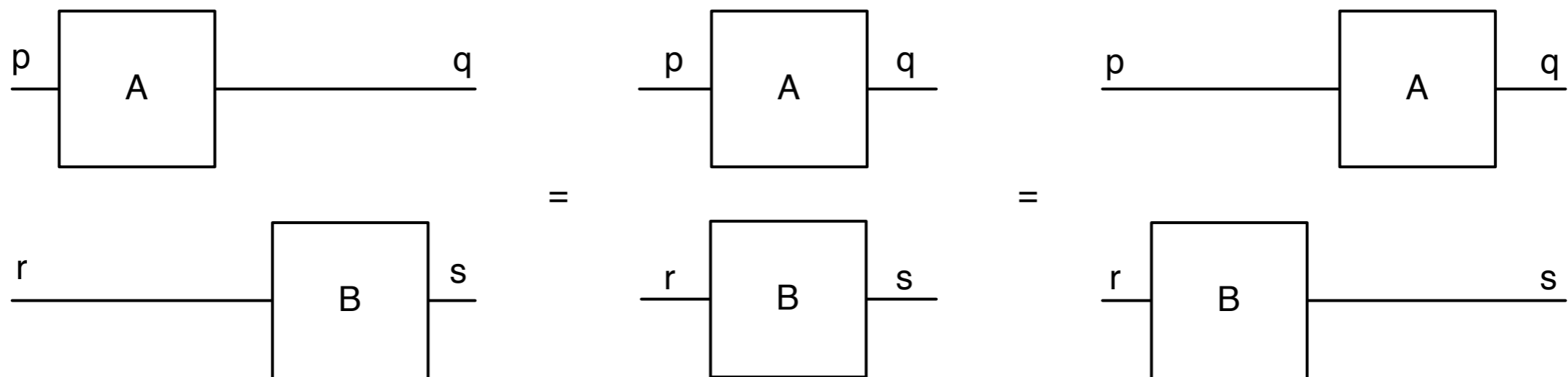


Diagrammatic reasoning

$$\frac{C : m \rightarrow n}{I_m ; C = C = C ; I_n}$$

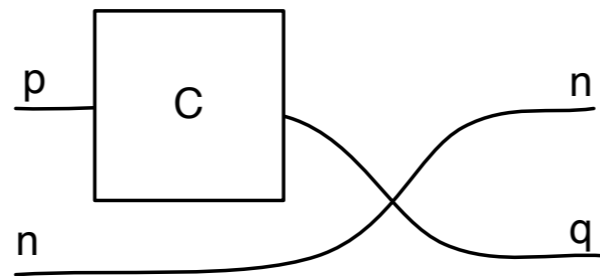
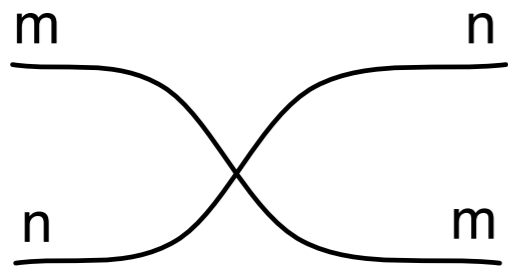


$$(A \oplus I_r) ; (I_q \oplus B) = A \oplus B = (I_p \oplus B) ; (A \oplus I_s)$$

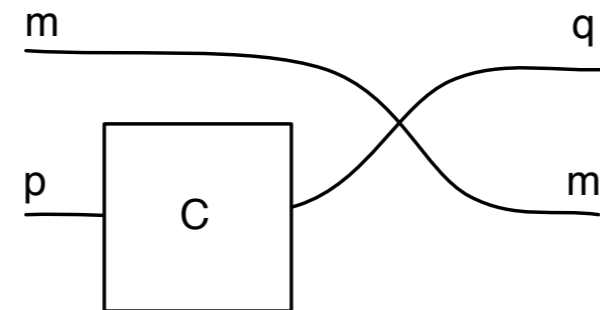
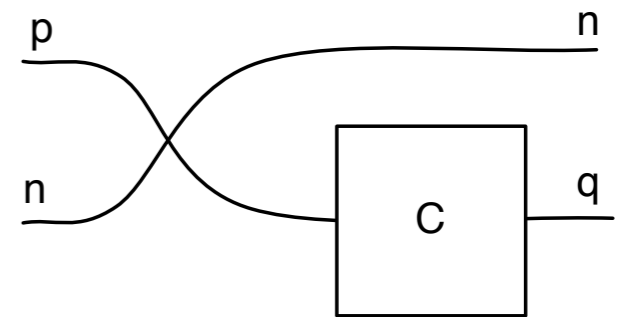


Symmetries

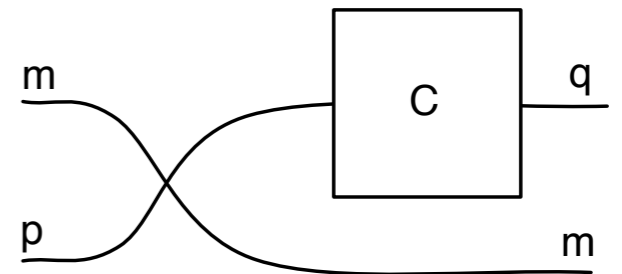
$$\sigma_{m,n}: m \oplus n \rightarrow n \oplus m$$



=



=



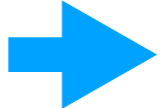
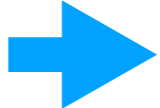
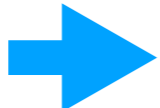
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(commutative) monoids and groups a la 1930s universal algebra - syntax

- (presentation of) algebraic theory
 - pair $T = (\Sigma, E)$ of finite sets
- for commutative monoids:
 - signature Σ , arity: $\Sigma \rightarrow \mathbf{N}$
 - $\cdot : 2$
 - $e : 0$
 - equations E (pairs of **typed terms**)
 - $x \cdot (y \cdot z) = (x \cdot y) \cdot z$
 - $x \cdot y = y \cdot x$
 - $x \cdot e = x$
- For abelian groups, additionally
 - signature: $(-)^{-1} : 1$
 - equations: $x \cdot x^{-1} = e$

(commutative) monoids and groups a la universal algebra - semantics

- To give a **model**
 - Pick carrier set X
 - $\cdot : 2 \rightarrow X$  $\cdot : X^2 \rightarrow X$
 - $e : 0 \rightarrow X$  $e : X^0 \rightarrow X$
 - $(-)^{-1} : 1 \rightarrow X$  $(-)^{-1} : X^1 \rightarrow X$
 - For every evaluation of variables $\sigma: \text{Var} \rightarrow X$, each equation must hold
- So, e.g. a **model of the algebraic theory of monoids** is the same thing as a **monoid**, in the classical sense

functorial semantics, 1960s

- Lawvere was not happy with universal algebra
 - too set theory specific
 - (e.g. topological groups morally should be a model)
 - too much ad hoc extraneous machinery
 - (e.g. countable set of variables, variable evaluation, etc.)
- Lawvere's 1963 doctoral thesis "Functorial semantics of algebraic theories" - universal algebra categorically

Lawvere theories

- Given algebraic theory (Σ, E) , a category $\mathcal{L}_{(\Sigma, E)}$ with

- objects: the natural numbers

- arrows from m to n :

- n -tuples of terms that (possibly) use variables x_1, x_2, \dots, x_m modulo equations E

- composition is substitution

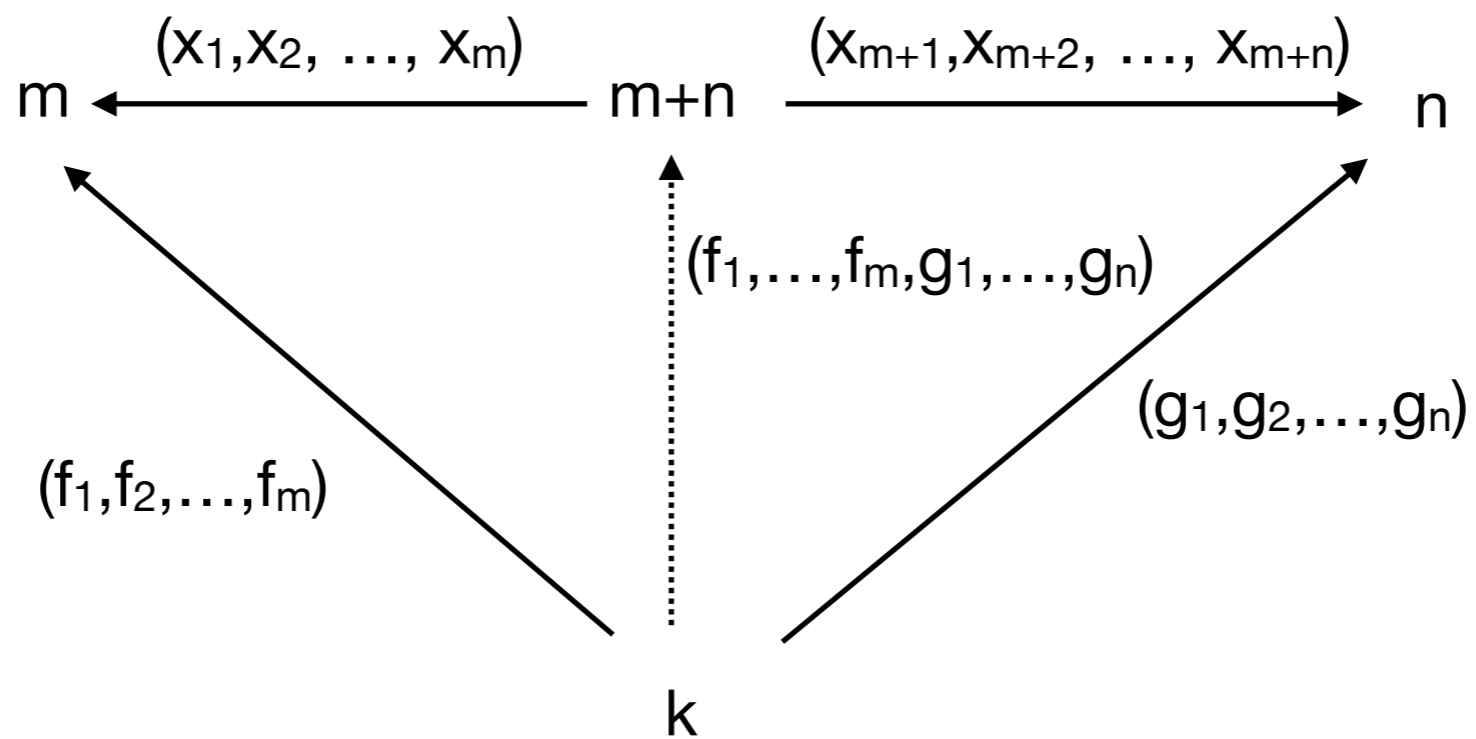
- e.g. $2 \xrightarrow{(x_1 \cdot x_2)} 1 \quad 2 \xrightarrow{(x_2 \cdot x_1)} 1 \quad 1 \xrightarrow{(e, x_1)} 2 \xrightarrow{(x_2 \cdot x_1)} 1 = 1 \xrightarrow{(x_1 \cdot e)} 1$

- More concisely - “free category with products on the data of an algebraic theory”

- any $\mathcal{L}_{(\Sigma, E)}$ is a prop!

classical model = cartesian functor $\mathcal{L} \rightarrow \mathbf{Set}$

products in a Lawvere theory



limitations of algebraic theories

- Copying and discarding **built in**

$$2 \xrightarrow{(X_1)} 1 \quad 2 \xrightarrow{(X_2)} 1 \quad 1 \xrightarrow{(X_1, X_1)} 2$$

- Consequently, there are also no bona fide operations with *coarities* other than one

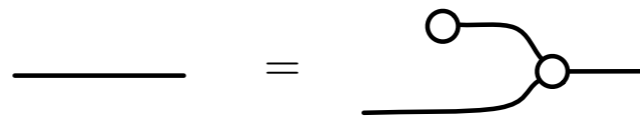
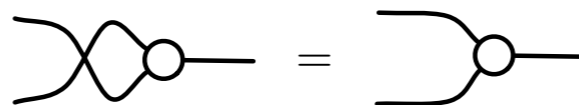
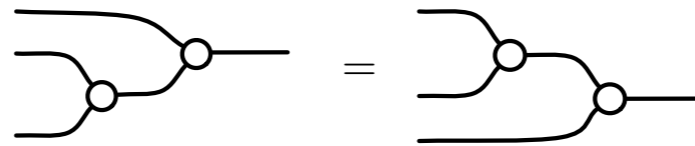
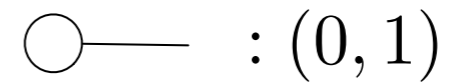
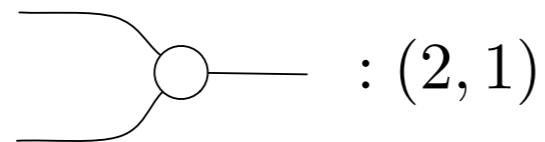
$$1 \xrightarrow{c} 2 \quad = \quad 1 \xrightarrow{(c_1, c_2)} 2$$

- But in quantum mechanics, computer science, and elsewhere we often need to be more careful with resources

symmetric monoidal theories

- algebraic theory in the symmetric monoidal settings
- a symmetric monoidal theory is a pair of finite sets (Σ, E)
 - Σ signature, arity : $\Sigma \rightarrow \mathbb{N}$, coarity : $\Sigma \rightarrow \mathbb{N}$
 - E equations, pairs of **string diagrams** constructed from Σ , identity and symmetries

symmetric monoidal theory of commutative monoids

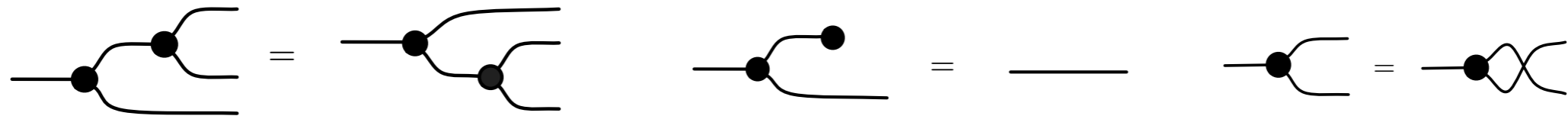


commutative monoid facts

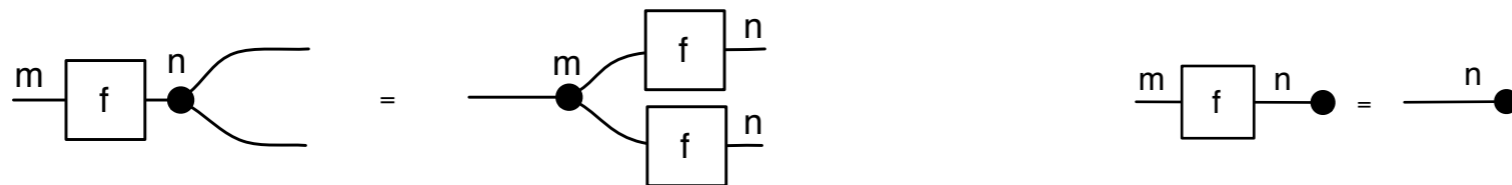
- the following are isomorphic as props
 - prop of commutative monoids
 - prop of functions between finite sets
- **not** isomorphic to the Lawvere theory of commutative monoids

folk theorem

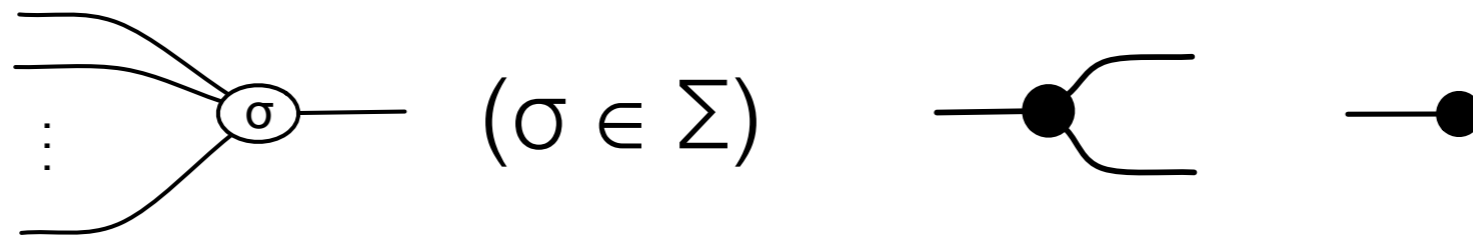
- A symmetric monoidal category \mathbf{C} is cartesian iff
 - every object $C \in \mathbf{C}$ has a commutative comonoid $\Delta: C \rightarrow C \otimes C$,
 $c: C \rightarrow I$



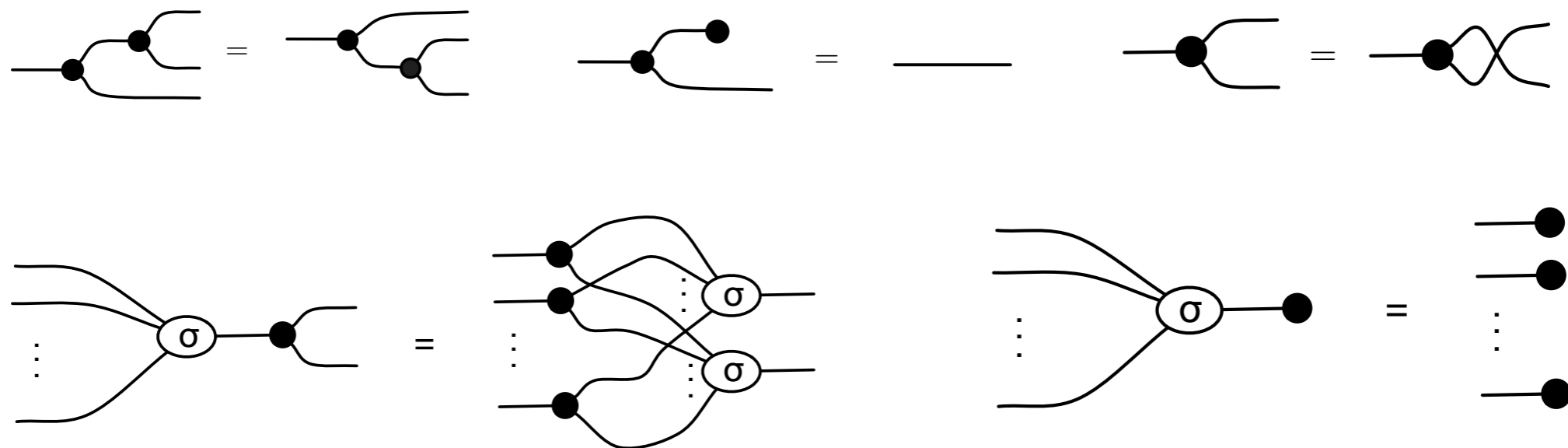
- compatible with \otimes in the obvious way
- and every arrow $f: m \rightarrow n$ of \mathbf{C} is a comonoid homomorphism, i.e.



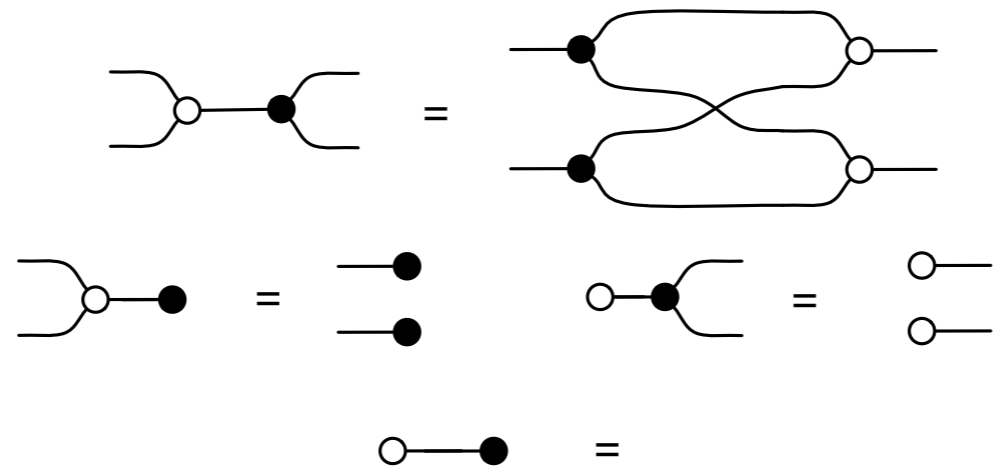
Lawvere theories as SMTs



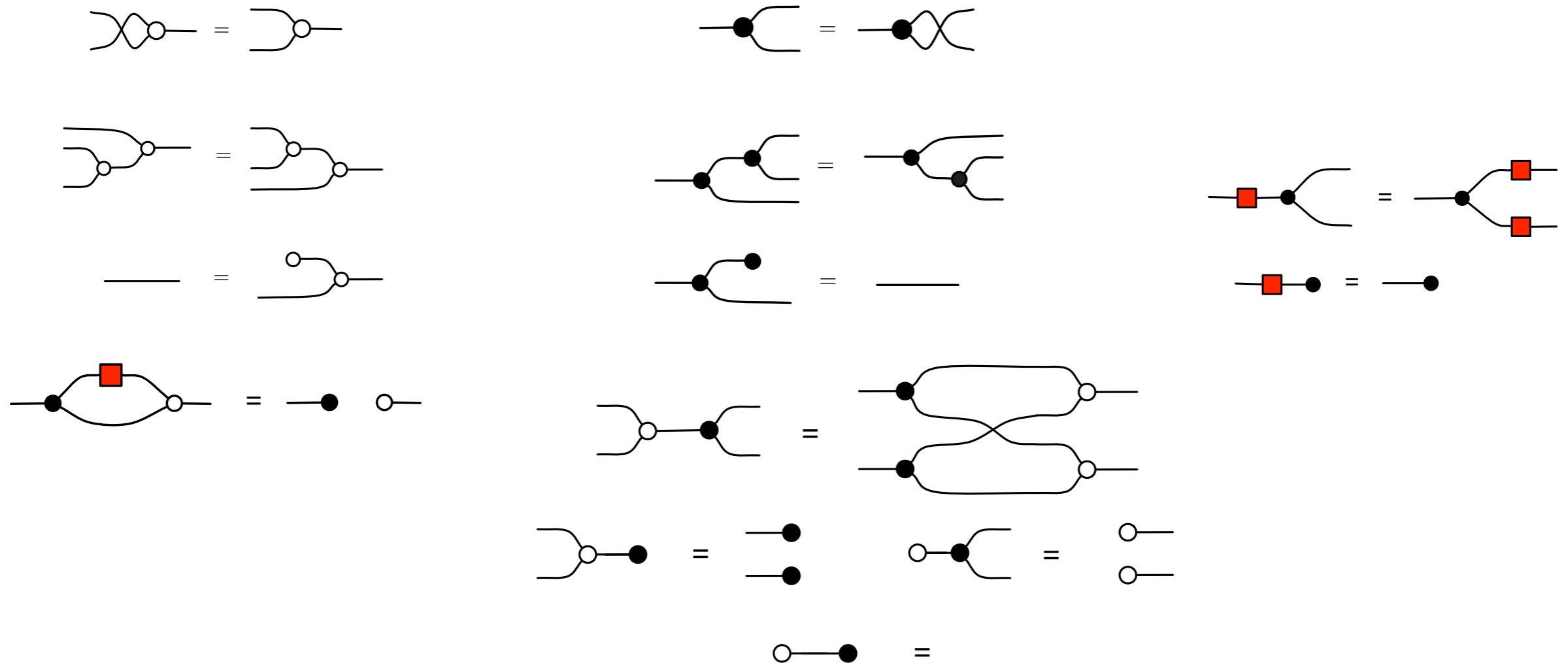
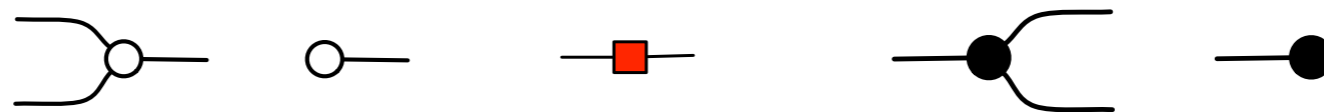
E +



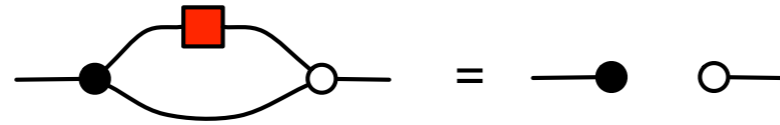
Lawvere theory of commutative monoids as SMT



Lawvere theory of abelian groups as an SMT



- e.g. the Hopf equation



is simply the SMT version of $x \cdot x^{-1} = e$

- Lawvere theory of commutative monoids = Symmetric monoidal theory of (co)commutative bialgebra
- Lawvere theory of abelian groups = Symmetric monoidal theory of (co)commutative Hopf algebras

So bialgebras and Hopf algebras are, respectively, monoids and groups in a resource sensitive universe.

Plan

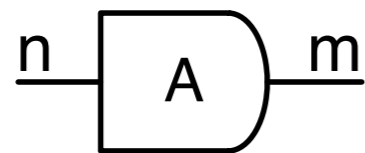
- String diagrams & diagrammatic reasoning
 - what is it?
 - why is it relevant for cs?
- **Graphical Linear Algebra**
- Fun stuff

Linear relation

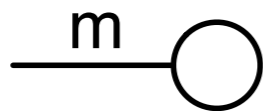
- Definition. Suppose V, W are k -vector spaces. A linear relation R from V to W is a linear subspace of $V \times W$
 - i.e.
 - $(0_V, 0_W) \in R$
 - if $(v, w), (v', w') \in R$ then $(v+v', w+w') \in R$
 - if $(v, w) \in R$ and $\lambda \in k$ then $(\lambda v, \lambda w) \in R$

Why linear relations?

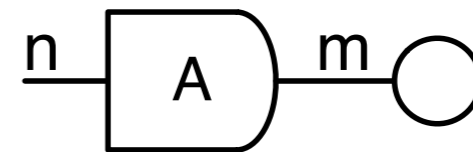
- any $m \times n$ matrix A gives lin. relation $\{ (\mathbf{x}, A\mathbf{x}) \mid \mathbf{x} \in k^n \} \subseteq k^n \times k^m$



- the singleton $(\mathbf{0}, *)$ is a linear relation $\subseteq k^m \times k^0$



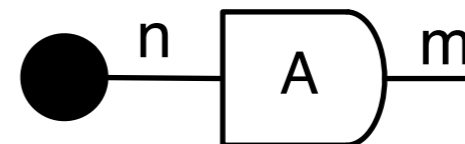
- composing gives the kernel of A



- the set $\{ (*, \mathbf{x}) \mid \mathbf{x} \in k^n \}$ is a linear relation $\subseteq k^0 \times k^n$



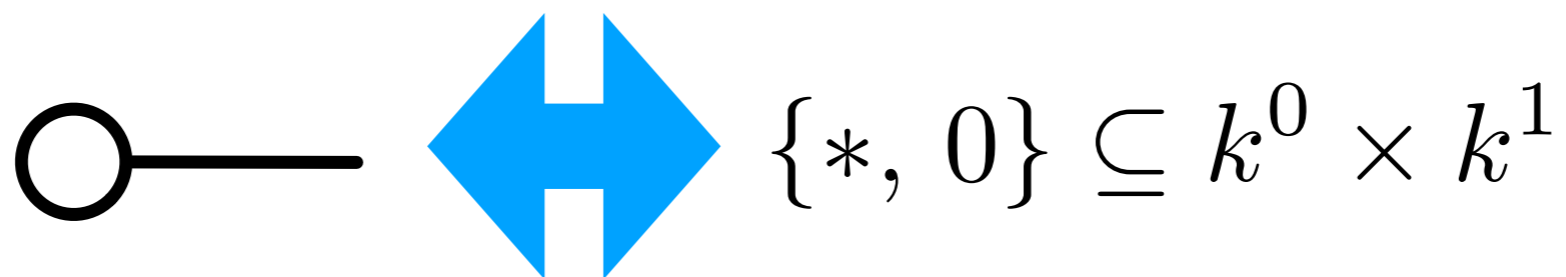
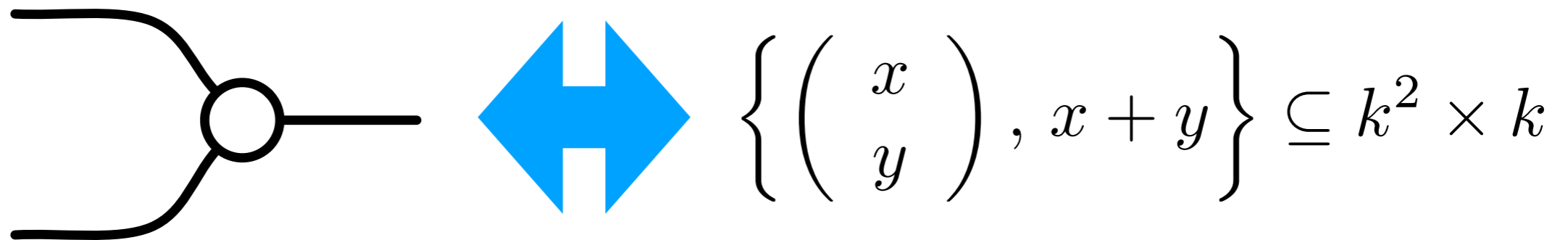
- composing gives the image of A



Graphical linear algebra

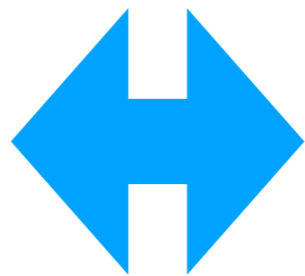
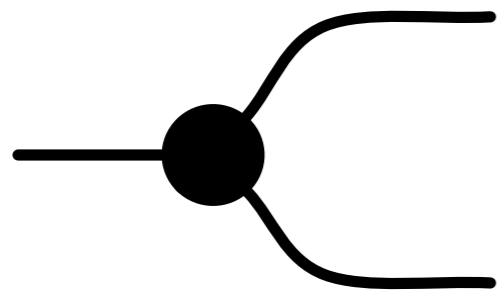
String diagrammatic syntax for linear relations with
a sound and fully complete axiomatisation called
Interacting Hopf Algebra

The signature, pt 1

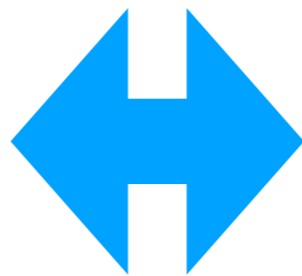
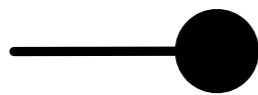


+ mirror images

The signature, pt 2



$$\left\{ x, \begin{pmatrix} x \\ x \end{pmatrix} \right\} \subseteq k \times k^2$$

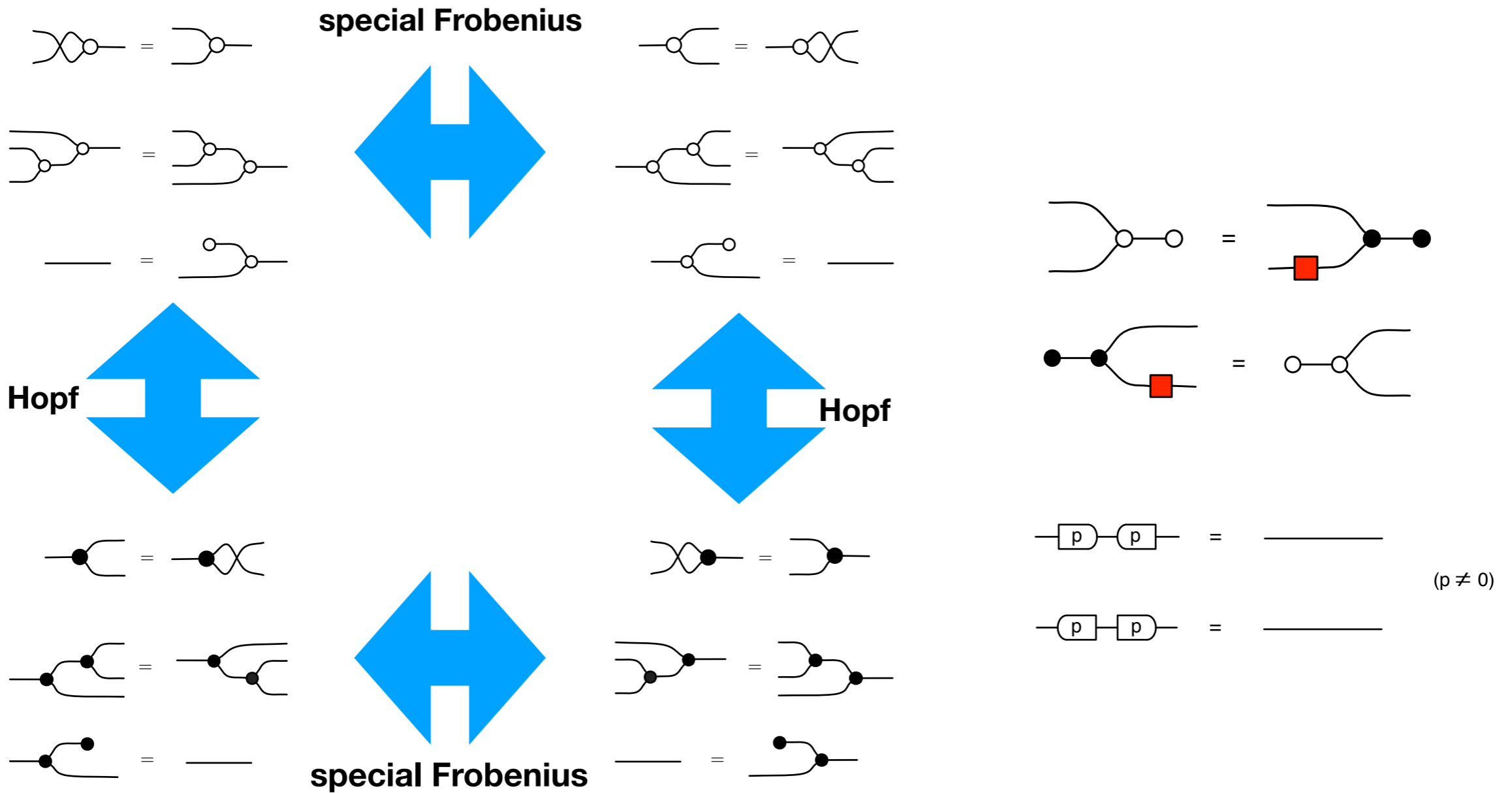
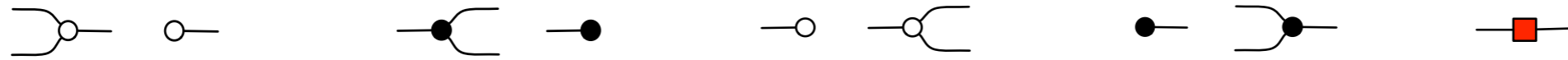


$$\{x, *\} \subseteq k \times k^0$$

+ mirror images

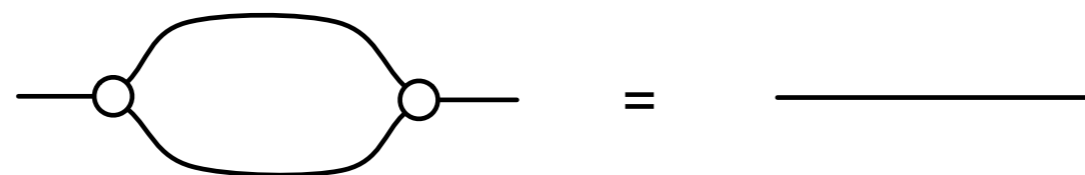
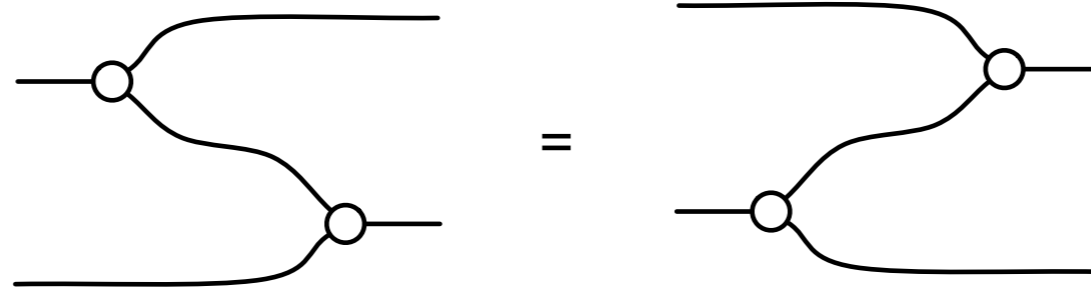
interacting Hopf algebras

Bonchi, S., Zanasi, JPAA 2017



cf. Coecke, Duncan. Interacting quantum observables, NJP 2011

(special) Frobenius monoids



Theorem

$$\text{IH} \cong \text{LinRel}$$



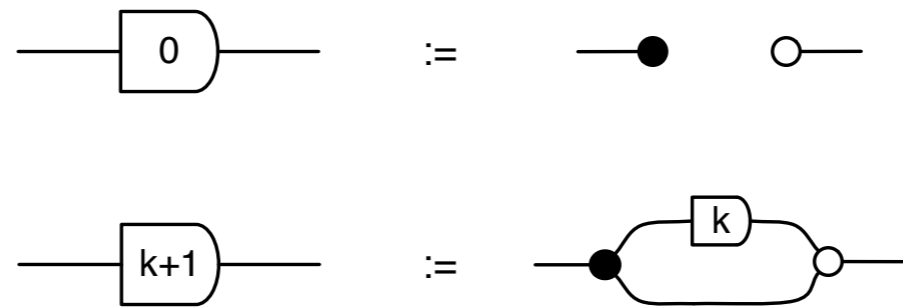
extends this to an iso of 2-categories

Plan

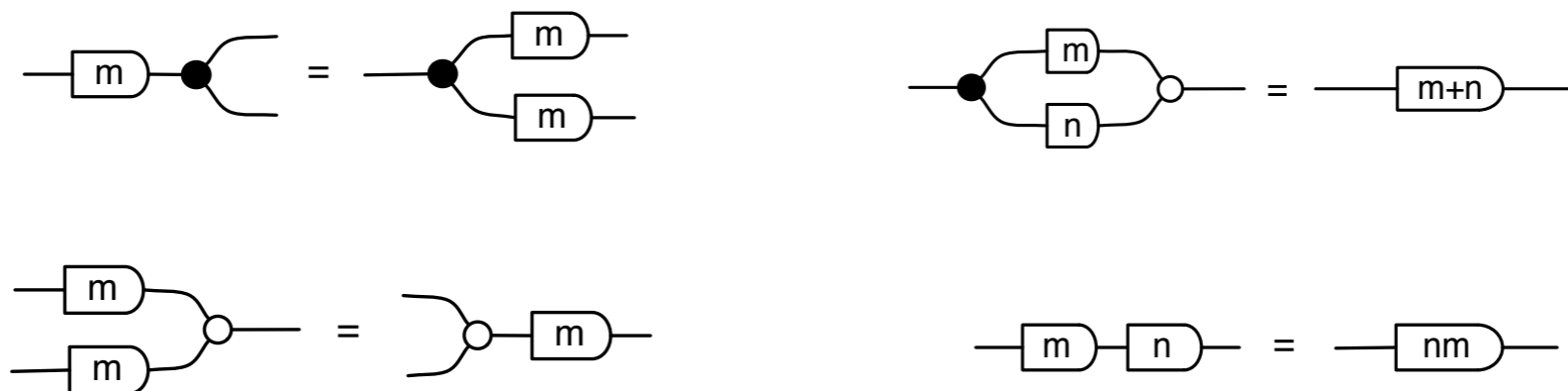
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naturals as string diagrams

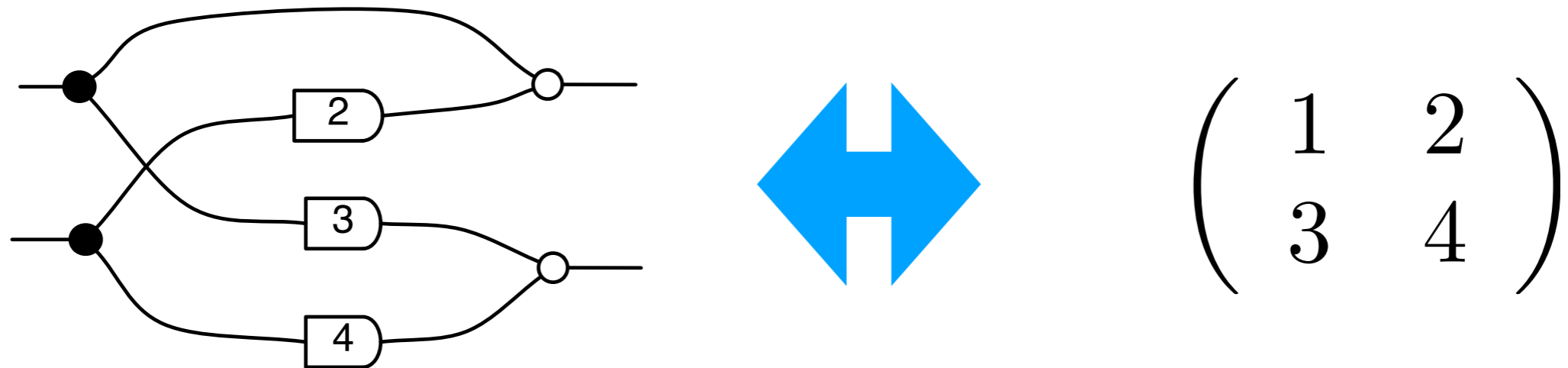
- naturals as *syntactic sugar*



- some easy lemmas

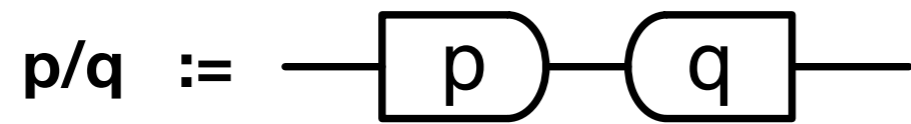


correspondence with matrices

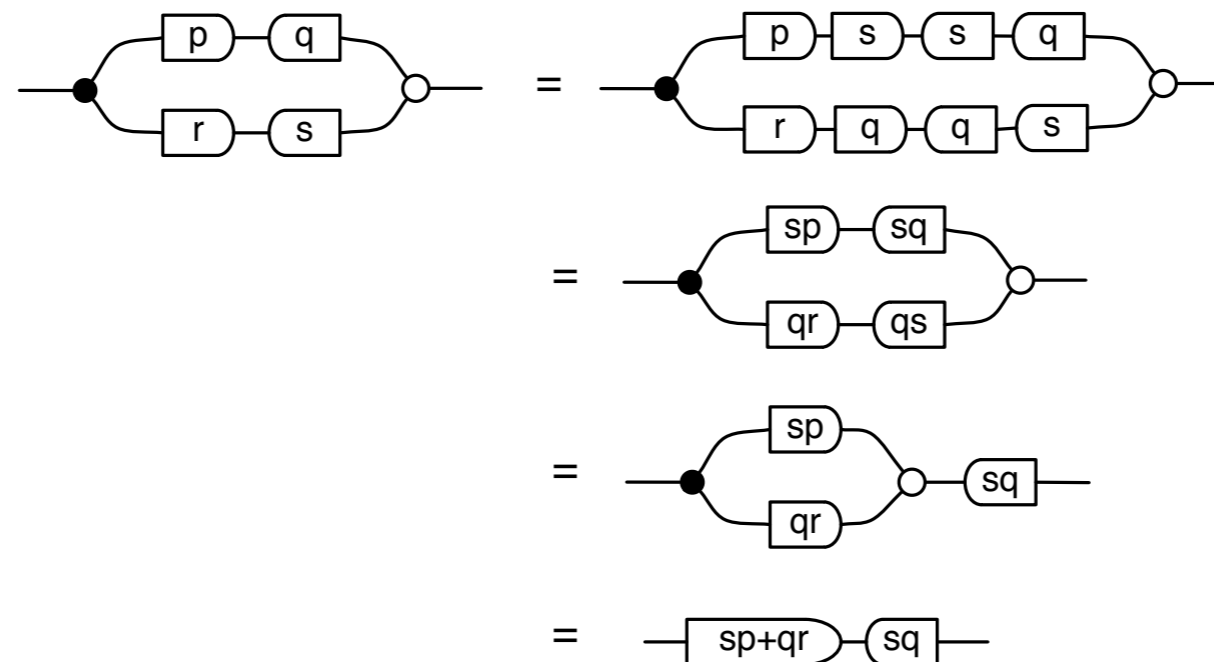
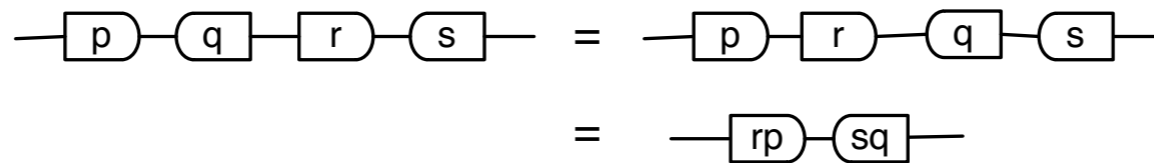
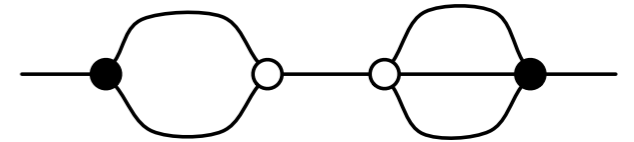


- in general, the ij th entry is the *number of paths* from the j th port on the left to the i th port on the right

rational numbers

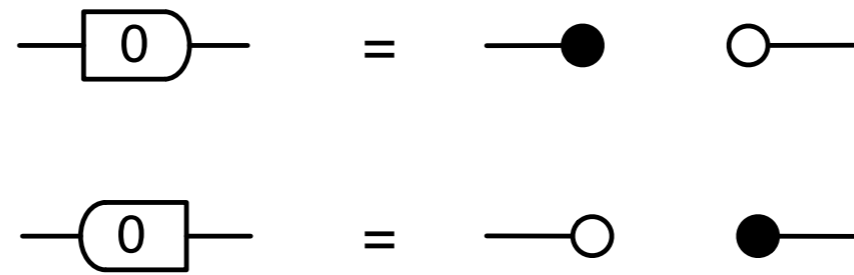


e.g. $2/3$ is

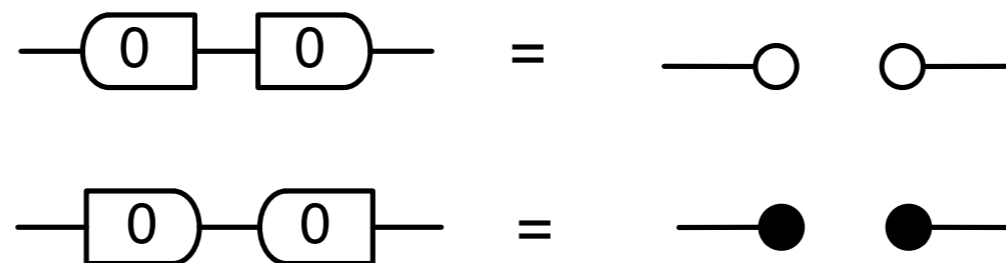


division by 0

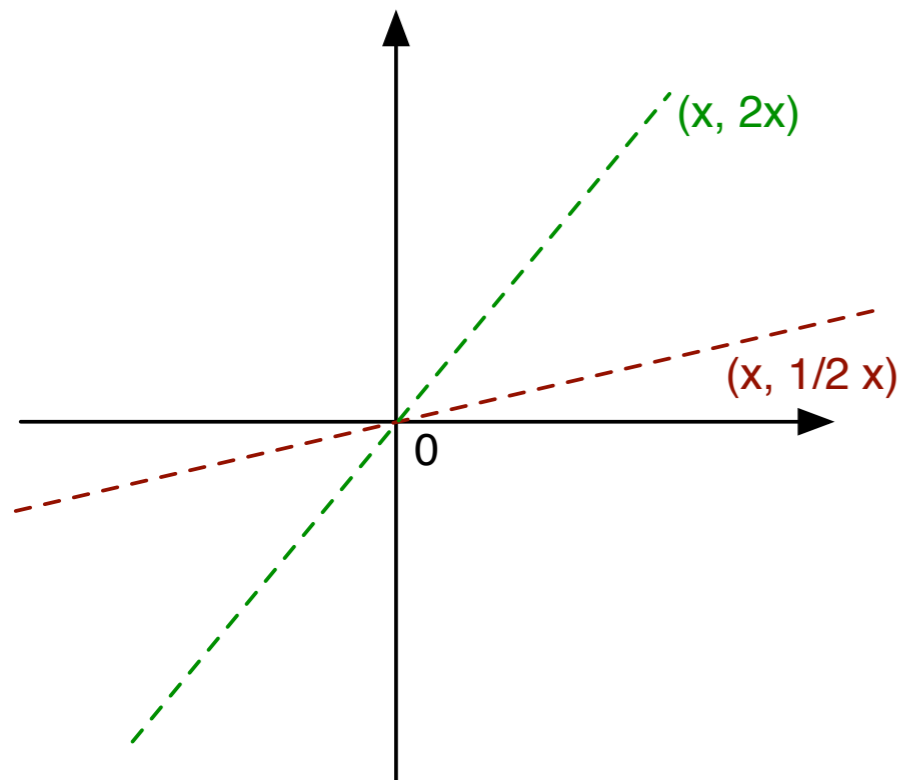
(fixing the deficiencies of the usual syntax)



Two ways of interpreting “0/0”



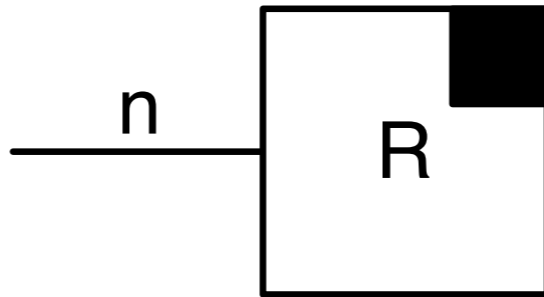
projective arithmetic ++



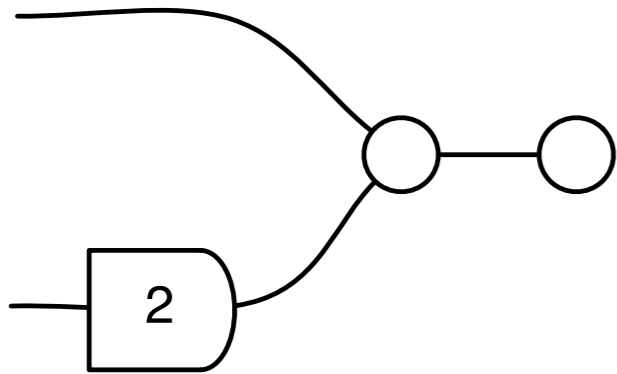
- projective arithmetic identifies rationals with 1-dim spaces (lines) of \mathbf{Q}^2
 - $p \rightarrow \{ (x, px) \mid x \in \mathbf{Q} \}$
 - $\infty : \{ (0, x) \mid x \in \mathbf{Q} \}$
- The extended system includes all the subspaces of \mathbf{Q}^2 , in particular:
 - the unique zero dimensional space $\{ (0, 0) \}$
 - the unique two dimensional space $\{ (x, y) \mid x, y \in \mathbf{Q} \}$

Linear subspaces

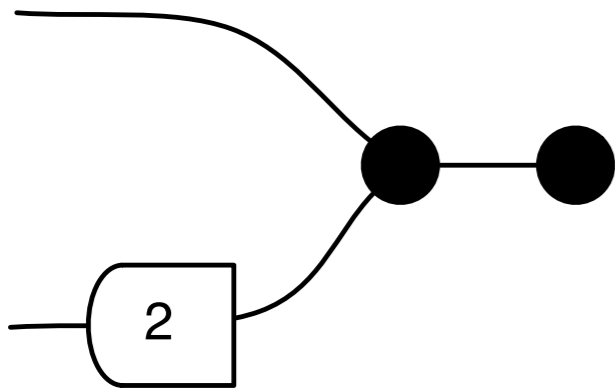
- **Observation.** Linear subspaces of \mathbb{R}^n are in 1-1 correspondence with string diagrams



Some examples

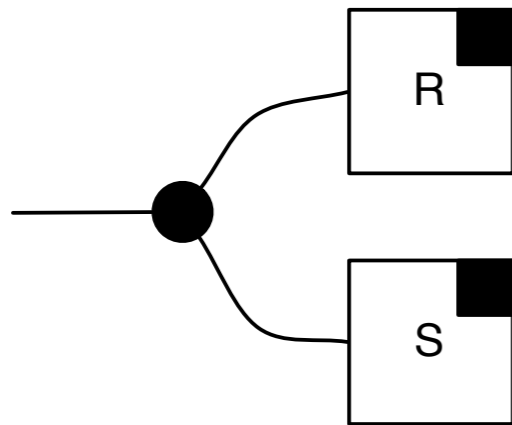


$$\left\{ \left(\begin{pmatrix} x \\ y \end{pmatrix}, * \right) \mid x + 2y = 0 \right\} \subseteq k^2 \times k^0$$

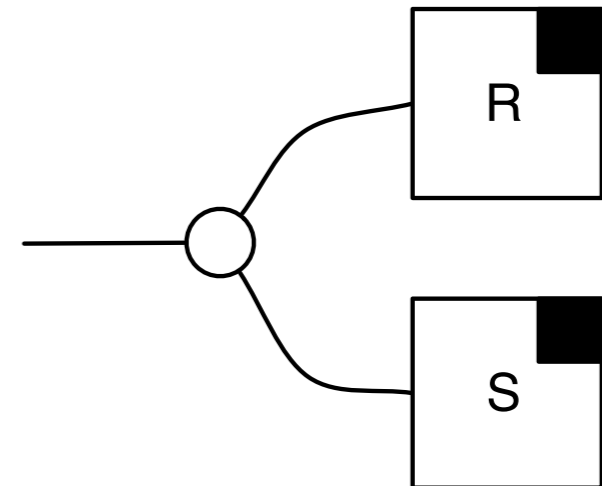


$$\left\{ \left(a \begin{pmatrix} 1 \\ 2 \end{pmatrix}, * \right) \mid a \in k \right\} \subseteq k^2 \times k^0$$

Intersection and sum of spaces

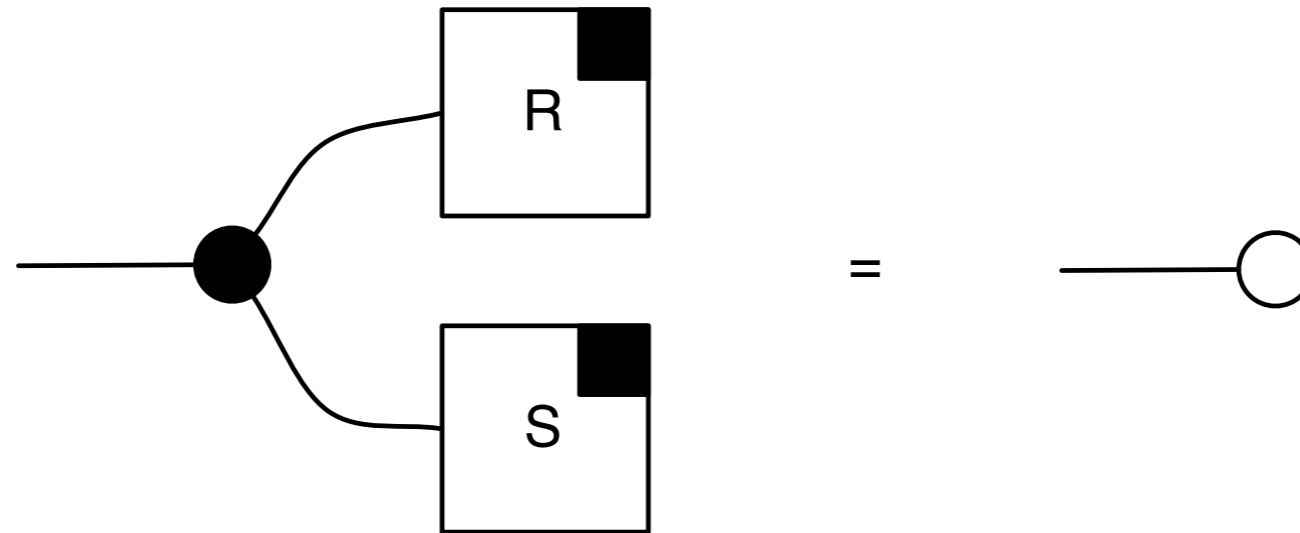


intersection

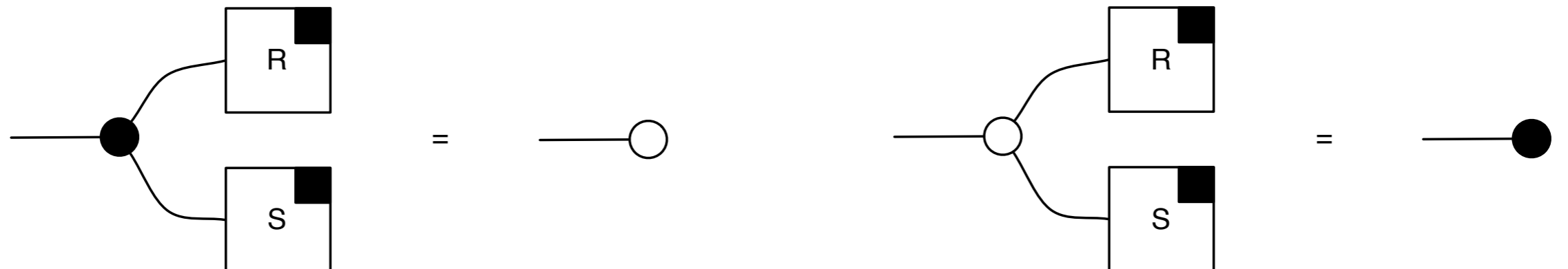


sum

linear independence

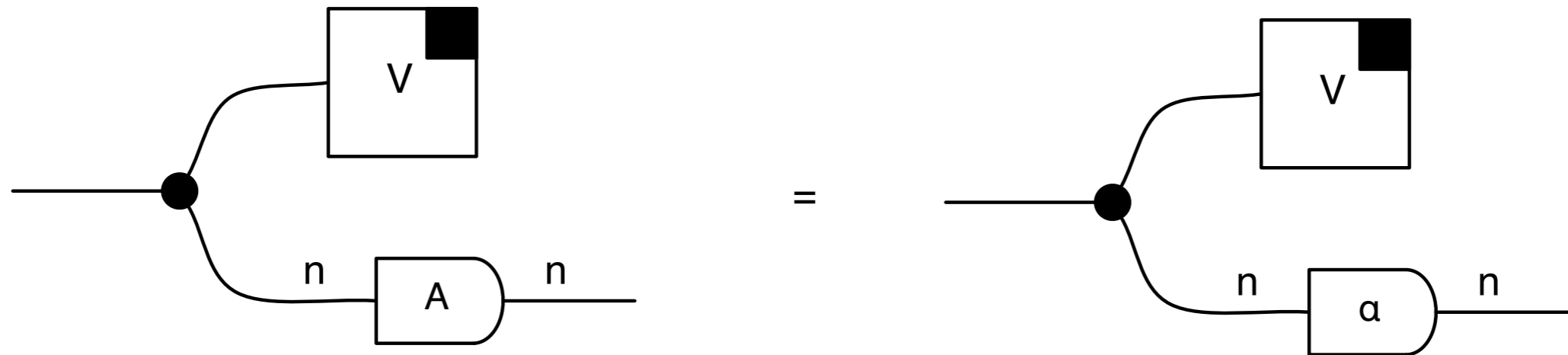


decomposition into linearly independent subspaces R and S



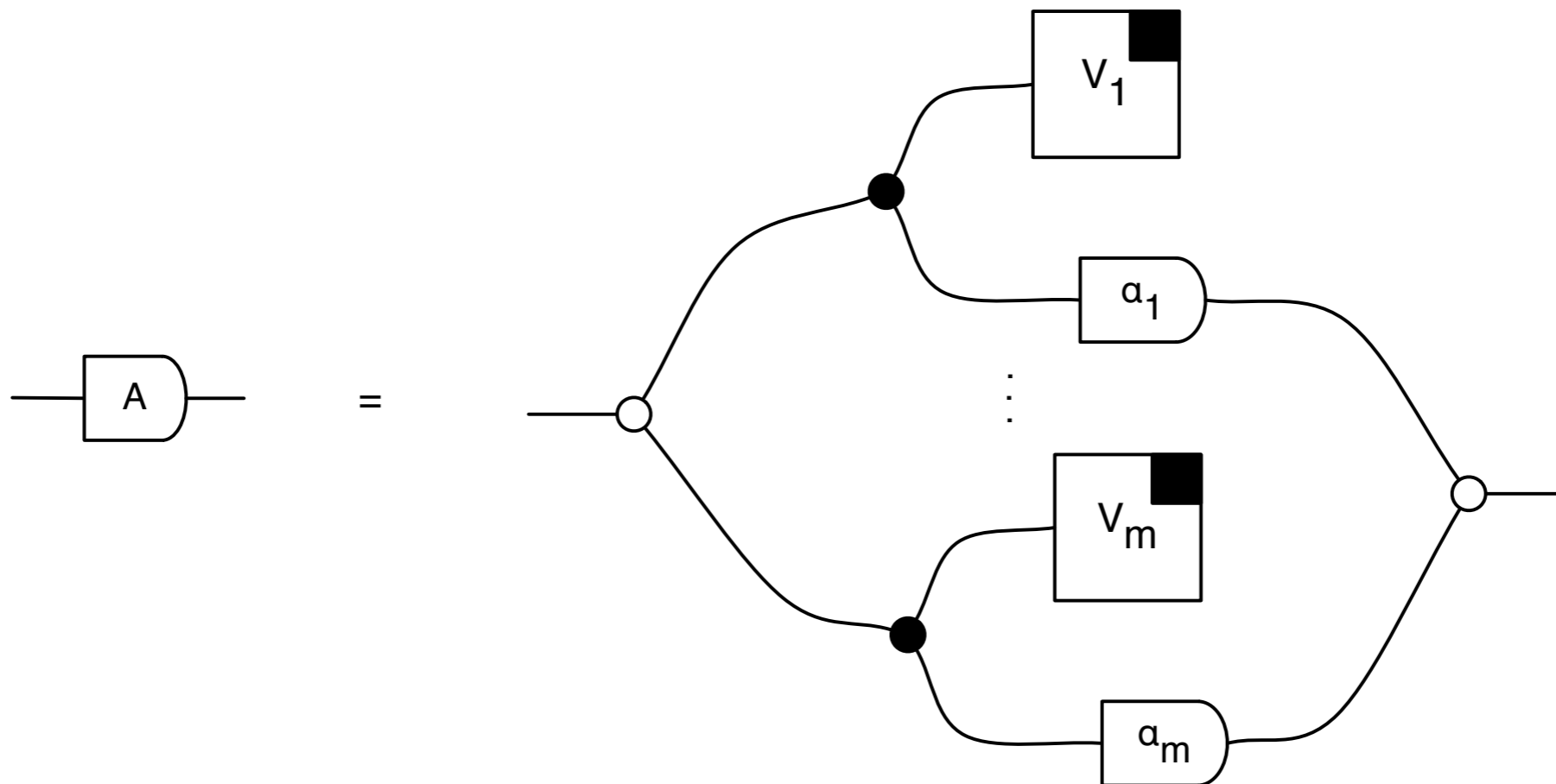
Eigenvalues & eigenspaces

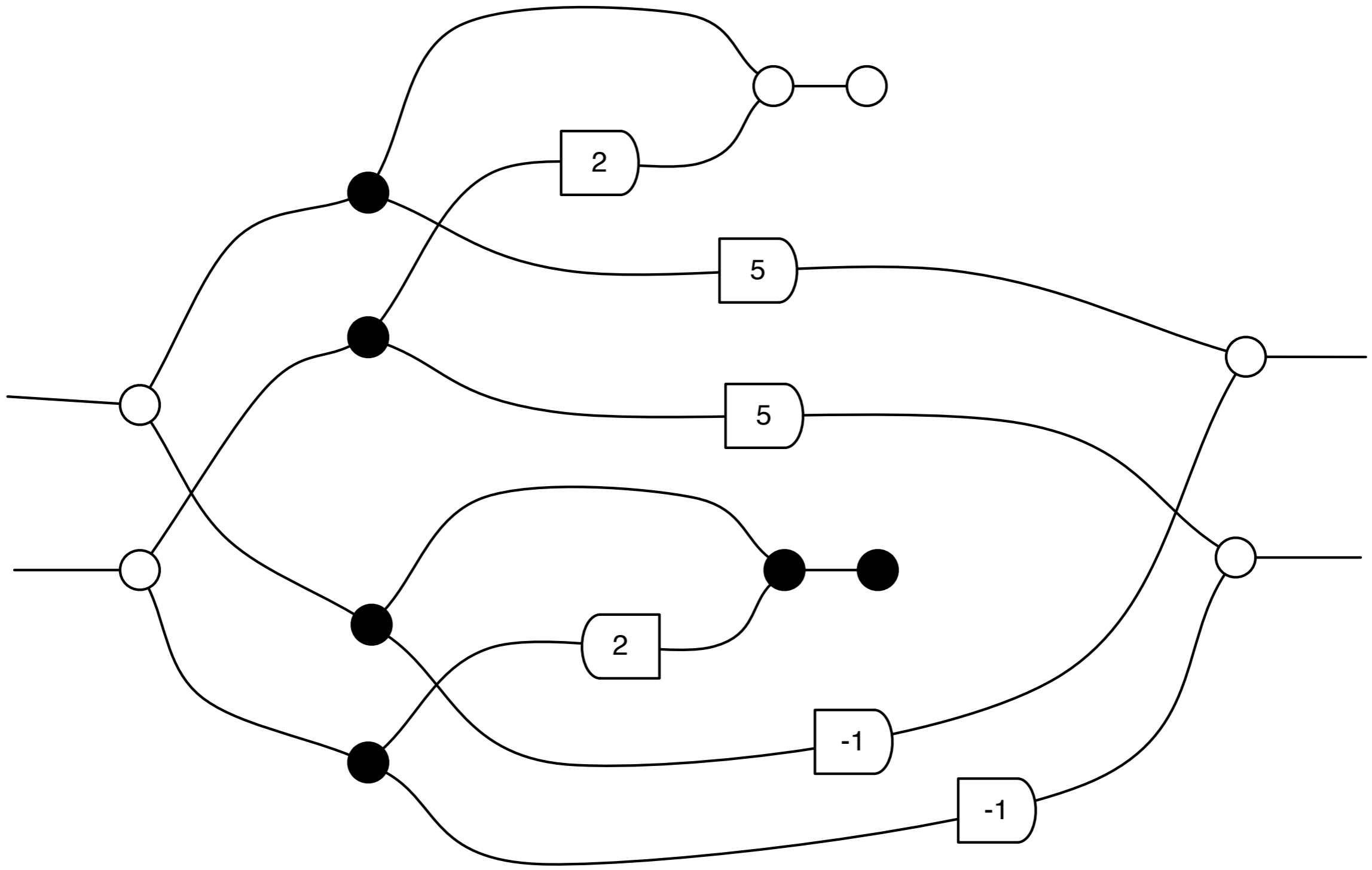
- V is an *eigenspace* of $A: k^n \rightarrow k^n$ with *eigenvalue* $a \in k$ when:

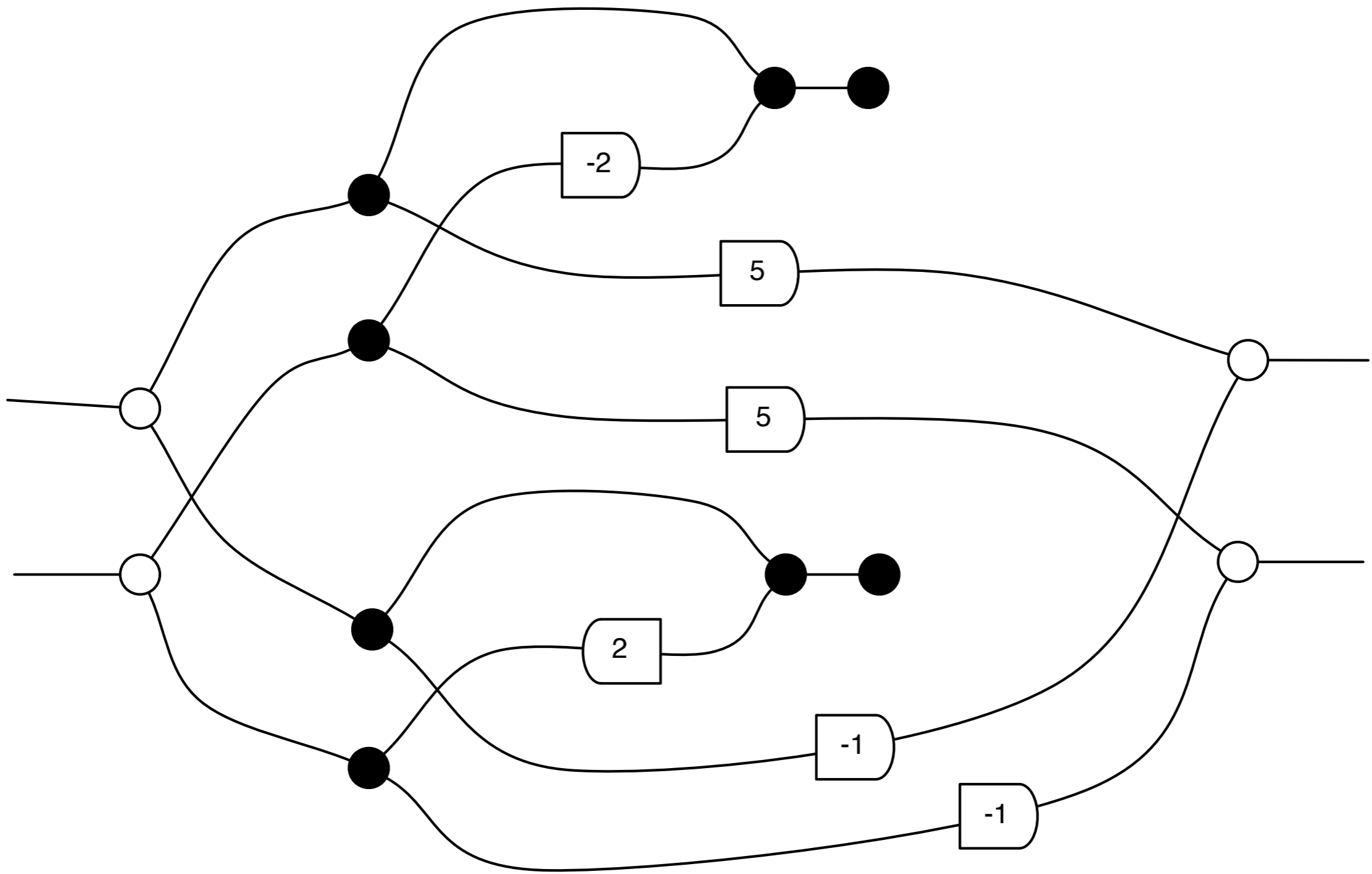


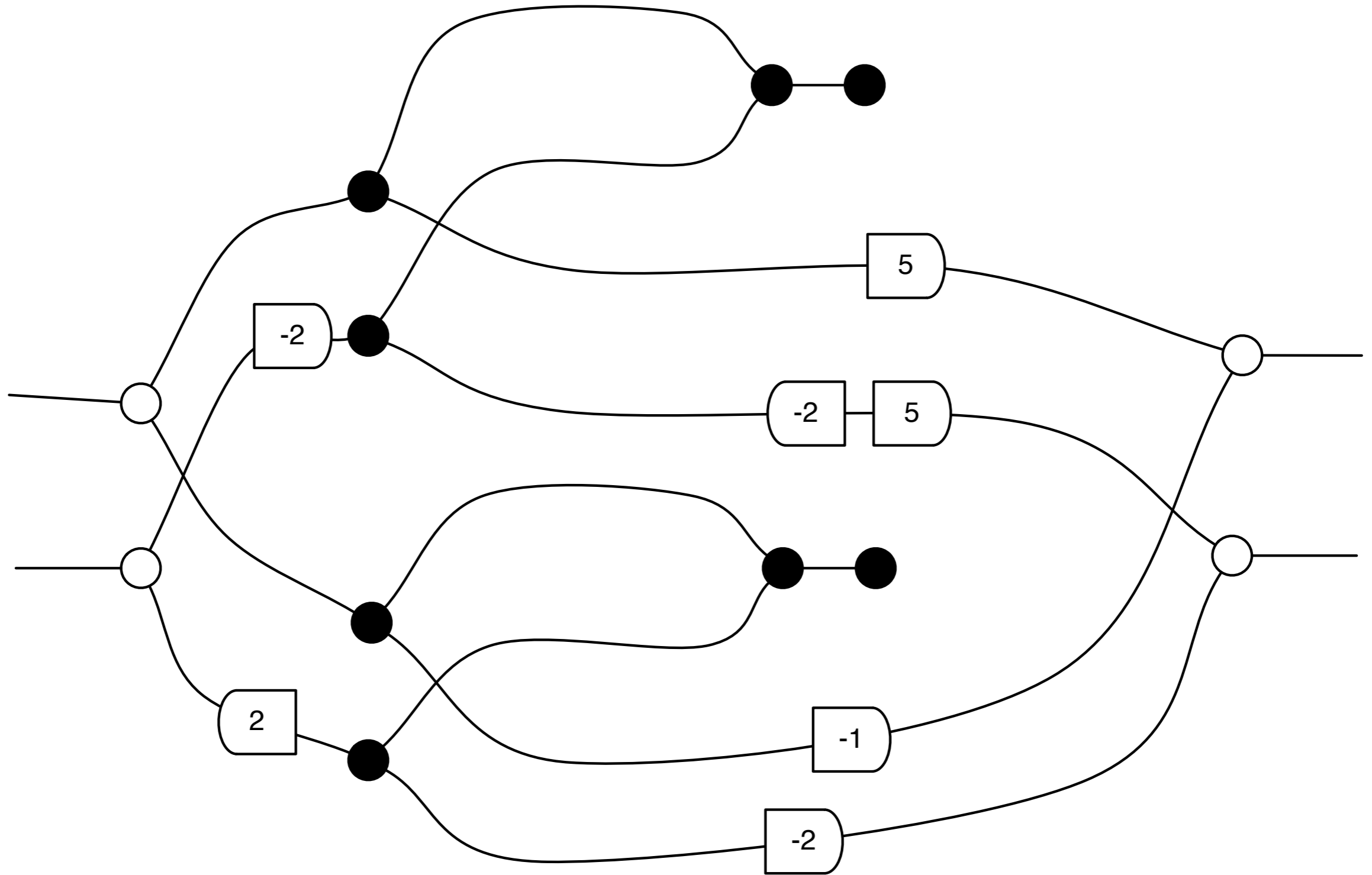
Spectral decomposition

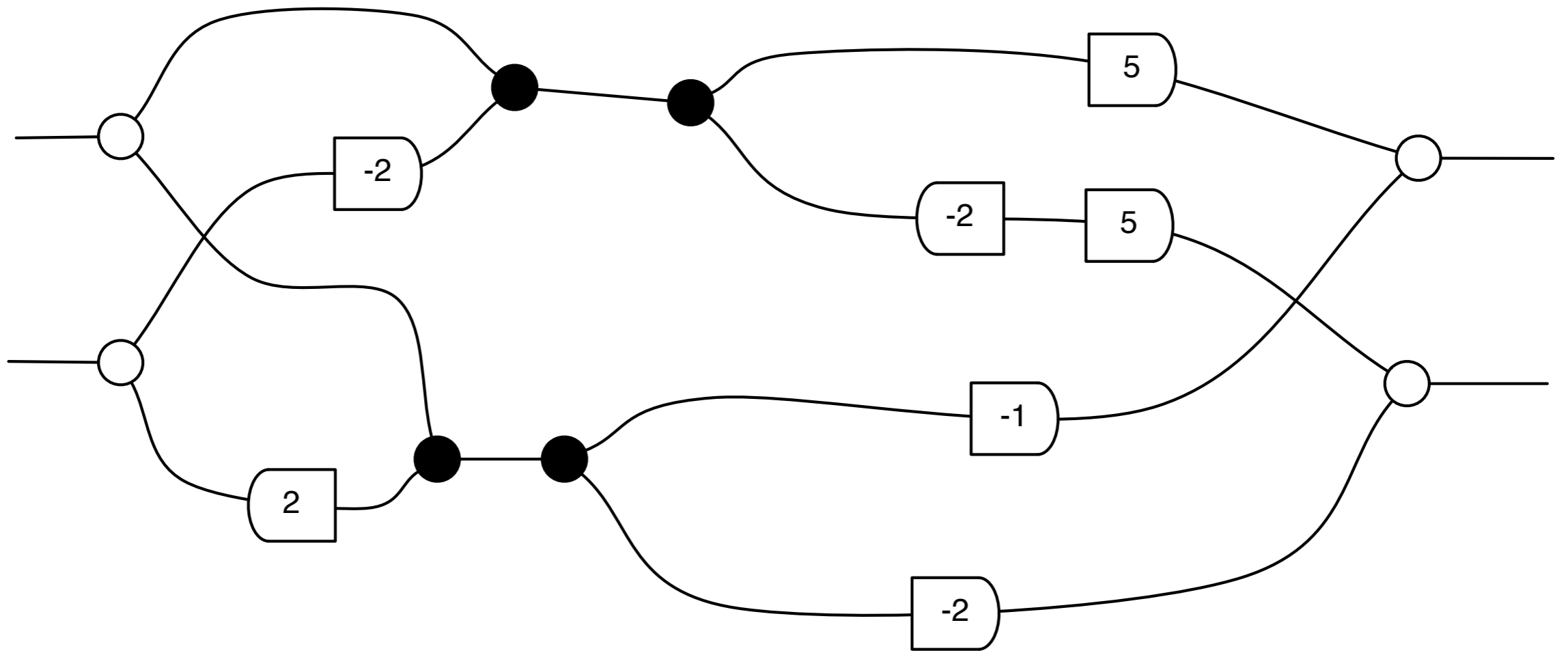
- A has a spectral decomposition when we can find a decomposition of k^n into eigenspaces V_1, V_2, \dots, V_m and eigenvalues $\alpha_1, \alpha_2, \dots, \alpha_m$











$$\begin{pmatrix} 5 & -1 \\ -\frac{1}{2} & -2 \end{pmatrix} \begin{pmatrix} 1 & 1 \\ -\frac{1}{2} & 2 \end{pmatrix}^{-1} = \frac{1}{5} \begin{pmatrix} 19 & -12 \\ -12 & 1 \end{pmatrix}$$

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